

Linear Codes with Two Weights or Three Weights from Two Types of Quadratic forms

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Abstract: Let \mathbb{F}_q be a finite field with $q = p^m$ elements, where p is an odd prime and m is a positive integer. Here, let $D_0 = \{(x_1, x_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\} : \text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) = c\}$, where $c \in \mathbb{F}_q$, Tr is the trace function from \mathbb{F}_q to \mathbb{F}_p and $m/(m, k_1)$ is odd, $m/(m, k_2)$ is even. Define a p -ary linear code $C_D = c(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2$, where $c(a_1, a_2) = (\text{Tr}(a_1 x_1 + a_2 x_2))_{(x_1, x_2) \in D}$. At most three-weight distributions of two classes of linear codes are settled.

Key words: linear codes; weight distribution; Gauss sum; Weil sum

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0 Introduction

Let \mathbb{F}_q be a finite field with $q = p^m$ elements throughout this paper, where p is an odd prime and m a positive integer, and let Tr be the trace function from \mathbb{F}_q to \mathbb{F}_p . An $[n, k, d]$ p -ary linear code C is a k -dimensional subspace of \mathbb{F}_q^n with the minimum Hamming distance d . Let A_i denote the number of codewords with Hamming weight i in code C of length n . The weight enumerator is defined by

$$1 + A_1 z + \cdots + A_n z^n$$

The sequence $(1, A_1, \dots, A_n)$ is called the weight distribution of the code C . A code C is said to be a t -weight code if the number of nonzero A_i is equal to t . The weight distribution is an interesting topic and was investigated in Refs. [1-9]. Weight distribution of code can not only give the error correcting ability of the code, but also allow the computation of the error probability of error detection and correction.

For a set $D = \{d_1, d_2, \dots, d_n\} \subseteq \mathbb{F}_q$, define a linear code of length n over \mathbb{F}_p by

$$C_D = \{(\text{Tr}(x d_1), \text{Tr}(x d_2), \dots, \text{Tr}(x d_n)) : x \in \mathbb{F}_q\}$$

We call D the defining set of C_D . Many known linear codes could be produced by the selected defining sets. For more details of these codes, please refer to Refs. [1, 3].

Here, assume that m, k_1, k_2 are the positive integers with $m/(m, k_1)$ being odd, $m/(m, k_2)$ being even. Then $f_1(x) = x^{p^{k_1}+1}$ is a planar function over $\mathbb{F}_q^{[10]}$. Fixing $c \in \mathbb{F}_p$, we define

$$D = \{(x_1, x_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\} : \\ \text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) = c\} \\ C_D = \{c(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2\}$$

where

$$c(a_1, a_2) = (\text{Tr}(a_1 x_1 + a_2 x_2))_{(x_1, x_2) \in D}$$

In fact, we have some well-known results as follows. If $a_1 = 0$, then it is just the result in Ref. [11].

We will determine the weight distribution of the linear codes C_D in three cases: (1) $c = 0$; (2) $c \in \mathbb{F}_p^{*2}$, (3) $c \in \mathbb{F}_p^* \setminus \mathbb{F}_p^{*2}$.

1 Preliminaries

Let \mathbb{F}_q be a finite field with q elements,

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where q is a power of a prime p . We define the additive character of \mathbb{F}_q as follows

$$\chi: \mathbb{F}_q \rightarrow \mathbb{C}^*, x \rightarrow \zeta_p^{Tr(x)}$$

where ζ_p is a complex p -th primitive root of unity and Tr the trace function from \mathbb{F}_q to \mathbb{F}_p . The orthogonal property of additive characters is given by

$$\sum_{x \in \mathbb{F}_q} \chi(ax) = \begin{cases} 0 & a \in \mathbb{F}_q^* \\ q & a = 0 \end{cases}$$

Let $\lambda: \mathbb{F}_q^* \rightarrow \mathbb{C}^*$ be a multiplicative character of \mathbb{F}_q^* . The trivial character λ_0 defined by $\lambda_0(x) = 1$ for all $x \in \mathbb{F}_q^*$. The orthogonal property of multiplicative characters is given by

$$\sum_{x \in \mathbb{F}_q^*} \lambda(x) = \begin{cases} q-1 & \lambda = \lambda_0 \\ 0 & \text{Otherwise} \end{cases}$$

Let $\bar{\lambda}$ be the conjugate character of λ and it is defined by $\bar{\lambda}(x) = \overline{\lambda(x)}$. It is easy to obtain that $\lambda^{-1} = \bar{\lambda}$. The multiplicative group \mathbb{F}_q^* is isomorphic to \mathbb{F}_q^* . For $\mathbb{F}_q^* = \langle \alpha \rangle$, define a multiplicative character by $\psi(\alpha) = \zeta_{q-1}$, where ζ_{q-1} denotes the primitive root of complex unity. Then we have $\mathbb{F}_q^* = \langle \psi \rangle$. Set $\eta = \psi^{\frac{q-1}{2}}$ the quadratic character of \mathbb{F}_q .

Define the Gauss sum over \mathbb{F}_q by

$$G(\lambda) = \sum_{x \in \mathbb{F}_q^*} \lambda(x) \chi(x)$$

Let $\left(\frac{\cdot}{p}\right)$ denote the Legendre symbol. The quadratic Gauss sums are known and given in the following.

Lemma 1.1^[12] Suppose that $q = p^m$ and η is the quadratic multiplicative character of \mathbb{F}_q , where p is odd prime. Then

$$G(\eta) = (-1)^{m-1} \sqrt{(p^*)^m} = \begin{cases} (-1)^{m-1} \sqrt{q} & p \equiv 1 \pmod{4} \\ (-1)^{m-1} (\sqrt{-1})^m \sqrt{q} & p \equiv 3 \pmod{4} \end{cases}$$

where $p^* = (-1)^{\frac{p-1}{2}} p$ is the discriminant of a prime p .

Lemma 1.2^[12] Let χ be a nontrivial additive character of the \mathbb{F}_q with q is odd, and let $f(x) = a_2 x^2 + a_1 x + a_0 \in \mathbb{F}_q[x]$ with $a_2 \neq 0$. Then

$$\sum_{c \in \mathbb{F}_q} \chi(f(c)) = \chi(a_0 - a_1^2 (4a_2)^{-1}) \eta(a_2) G(\eta, \chi)$$

where η is the quadratic character of \mathbb{F}_q .

Let χ' be the canonical additive character of \mathbb{F}_p such that $\chi(x) = \chi'(Tr(x))$ for $x \in \mathbb{F}_q$. Let η' be a quadratic character of \mathbb{F}_p , then $\eta(x) = \eta'(N_{q/p}(x))$ for $x \in \mathbb{F}_q^*$.

Lemma 1.3^[13] Let $x \in \mathbb{F}_p^*$ and $q = p^m$, where p is odd prime.

If m is even, $\eta(x) = 1$.

If m is odd, $\eta(x) = \eta'(x)$.

Moreover, $G(\eta) = (-1)^{m-1} G(\eta')^m$, where $G(\eta)$ and $G(\eta')$ are the Gauss sums over \mathbb{F}_q and \mathbb{F}_p , respectively.

We now give a brief introduction to the theory of quadratic forms over finite fields. Quadratic forms have been well studied and have been applied to sequence design^[9,14] and coding theory^[15].

Lemma 1.4 Let $d = \gcd(k, m)$. Then

$$(p^k + 1, p^m - 1) = \begin{cases} 2 & m/d \text{ is odd} \\ p^d + 1 & m/d \text{ is even} \end{cases}$$

In Refs. [16,17], Coulter gave the valuation of the following Weil sums

$$S_k(a, b) = \sum_{x \in \mathbb{F}_q} \chi(ax^{p^k+1} + bx) \quad a, b \in \mathbb{F}_q$$

Lemma 1.5^[16] Let m/d be odd. Then

$$S_k(a, 0) = \eta(a) G(\eta) = \begin{cases} (-1)^{m-1} \sqrt{q} \eta(a) & p \equiv 1 \pmod{4} \\ (-1)^{m-1} i^m \sqrt{q} \eta(a) & p \equiv 3 \pmod{4} \end{cases}$$

Lemma 1.6^[16] Let m/d be even with $m = 2e$. Then

$$S_k(a, 0) =$$

$$\begin{cases} p^e & a^{(q-1)/(p^d+1)} \neq (-1)^{e/d} \text{ and } e/d \text{ is even} \\ -p^e & a^{(q-1)/(p^d+1)} \neq (-1)^{e/d} \text{ and } e/d \text{ is odd} \\ p^{e+d} & a^{(q-1)/(p^d+1)} = (-1)^{e/d} \text{ and } e/d \text{ is odd} \\ -p^{e+d} & a^{(q-1)/(p^d+1)} = (-1)^{e/d} \text{ and } e/d \text{ is even.} \end{cases}$$

Lemma 1.7^[17] Let q be odd and suppose $f(X) = a^k X^{p^{2k}} + aX$ is a permutation polynomial over F_q . Let x_0 be the unique solution of the equation $f(X) = -b^{p^k}$. The evaluation of $S_k(a, b)$ partitions into the following two cases:

(1) If m/d is odd

$$S_k(a, b) = \eta(a) G(\eta) \bar{\chi}(ax_0^{p^k+1}) =$$

$$\begin{cases} (-1)^{m-1} \sqrt{q} \eta(a) \bar{\chi}(ax_0^{b^k+1}) & p \equiv 1 \pmod{4} \\ (-1)^{m-1} i^m \sqrt{q} \eta(a) \bar{\chi}(ax_0^{b^k+1}) & p \equiv 3 \pmod{4} \end{cases}$$

(2) If m/d is even, $m = 2e$, and $a_{p^d+1}^{\frac{q-1}{p^d}} \neq (-1)^{e/d}$

$$S_k(a, b) = (-1)^{e/d} p^e \bar{\chi}(ax_0^{b^k+1})$$

In fact, Lemma 1.4 is a revised version of the lemma in Ref. 4.

Lemma 1.8^[17] Let q be odd and suppose $f(X) = a^{b^k} X^{p^{2k}} + aX$ is not a permutation polynomial over F_q . Then for $b \neq 0$, we have $S_k(a, b) = 0$ unless the equation $f(X) = -b^{b^k}$ is solvable. If the equation is solvable, with some solution x_0 say, $S_k(a, b) = -(-1)^{e/d} p^{e+d} \bar{\chi}(ax_0^{b^k+1})$.

Lemma 1.9^[11] Let $f(X) = X^{p^{2k}} + X$ and $S = \{b \in \mathbb{F}_q : f(X) = -b^{b^k} \text{ is solvable in } \mathbb{F}_q\}$, If m/d is even, $|S| = p^{2d}$.

2 Linear codes

Let \mathbb{F}_q be the finite field with $q = p^m$ elements, where p is an odd prime and m an even positive integer with $m = 2e$. Let Tr denote the trace function from \mathbb{F}_q to \mathbb{F}_p . In this section, we always assume that e, d, k_1, k_2 are the positive integers with $m/\gcd(m, k_1)$ odd, $m/\gcd(m, k_2)$ even and $d = \gcd(m, k_2)$. Let $f_i(x) = x^{b^{k_i}+1}$, $x \in \mathbb{F}_q, i = 1, 2$.

2.1 The first case

Define

$$D_0 = \{(x_1, x_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\} : \text{Tr}(x_1^{b^{k_1}+1} + x_2^{b^{k_2}+1}) = 0\} \quad (1)$$

$$C_{D_0} = \{c(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2\}$$

where $c(a_1, a_2) = (Tr(a_1 x_1 + a_2 x_2))_{(x_1, x_2) \in D_0}$.

Lemma 2.1 Let $n_0 = |D_0|$. Suppose that e/d is odd, then

$$n_0 = \frac{q^2 - p}{p} + \frac{p-1}{p} (-p^e) G(\eta) = \begin{cases} p^{2m-1} - 1 + (-1)^m (p-1) p^{m-1} & p \equiv 1 \pmod{4} \\ p^{2m-1} - 1 + (-1)^{m+\frac{m}{2}} (p-1) p^{m-1} & p \equiv 3 \pmod{4} \end{cases}$$

Suppose that e/d is even, then

$$n_0 = \frac{q^2 - p}{p} + \frac{p-1}{p} (-p^{e+d}) G(\eta) =$$

$$\begin{cases} p^{2m-1} - 1 + (-1)^m (p-1) p^{m+d-1} & p \equiv 1 \pmod{4} \\ p^{2m-1} - 1 + (-1)^{m+\frac{m}{2}} (p-1) p^{m+d-1} & p \equiv 3 \pmod{4} \end{cases}$$

Proof 1 By Lemmas 1.4, 1.5 and 1.6, we have that

$$\begin{aligned} n_0 + 1 &= \frac{1}{p} \sum_{y \in \mathbb{F}_p} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \chi(y(x_1^{b^{k_1}+1} + x_2^{b^{k_2}+1})) = \\ &= \frac{q^2}{p} + \frac{1}{p} \sum_{y \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{b^{k_1}+1}) \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{b^{k_2}+1}) = \\ &= \begin{cases} \frac{q^n}{p} + \frac{1}{p} \sum_{y \in \mathbb{F}_p^*} G(\eta) \eta(y) (-p^e) & e/d \text{ is odd} \\ \frac{q^n}{p} + \frac{1}{p} \sum_{y \in \mathbb{F}_p^*} G(\eta) \eta(y) (-p^{e+d}) & e/d \text{ is even} \end{cases} \end{aligned}$$

$$\begin{cases} \frac{q^2}{p} + \frac{p-1}{p} (-p^e) G(\eta) & e/d \text{ is odd} \\ \frac{q^2}{p} + \frac{p-1}{p} (-p^{e+d}) G(\eta) & e/d \text{ is even} \end{cases}$$

For each $y \in \mathbb{F}_p^*$, if m is even, $2 \mid \frac{p^m-1}{p-1}$ and $\eta(y) = 1$;

For each $y \in \mathbb{F}_p^*$, we can know $y^{\frac{q-1}{p^d}} = 1$. Then by Lemma 1.5, we can obtain the exact value of n_0 .

Theorem 2.2 Let C_{D_0} be the linear code defined as Section 2.1.

If e/d is odd, C_{D_0} is a two-weight code with the Hamming weight distribution in Table 1.

If e/d is even, C_{D_0} is a three-weight code with the Hamming weight distribution in Table 2.

Table 1 e/d is odd

$p \equiv 1 \pmod{4}$	
Weight	Multiplicity
0	1
$p^{2m-1} - p^{2m-2}$	$\frac{q^2 - p}{p} + (p-1) p^{m-1}$
$(p-1)(p^{2m-2} + (-1)^m p^{m-1})$	$\frac{p-1}{p} (p^{2m} + (-1)^{m-1} p^m)$

$p \equiv 3 \pmod{4}$

Weight	Multiplicity
0	1
$p^{2m-1} - p^{2m-2}$	$\frac{q^2 - p}{p} + (-1)^{\frac{3m}{2}} (p-1) p^{m-1}$
$(p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}} p^{m-1})$	$\frac{p-1}{p} (p^{2m} + (-1)^{(m-1)+\frac{m}{2}} p^m)$

Table 2 e/d is even

$$p \equiv 1 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)(p^{2m-2} + (-1)^m \cdot (p-1)p^{m+d-2})$	$p^{2m} - p^{m+2d}$
$p^{2m-1} - p^{2m-2}$	$(p-1)p^{m-d-1} + p^{m+2d-1} - 1$
$(p-1)(p^{2m-2} + (-1)^m p^{m+d-1})$	$(p-1)(p^{m+2d-1} - p^{m-d-1})$

$$p \equiv 3 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}}(p-1)p^{m+d-2})$	$p^{2m} - p^{m+2d}$
$p^{2m-1} - p^{2m-2}$	$(p-1)p^{m-d-1} + p^{m+2d-1} - 1$
$(p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}}p^{m+d-1})$	$(p-1)(p^{m+2d-1} - p^{m-d-1})$

Proof 2 Firstly, we determine the weight distribution of the code C_{D_0} . Define the following parameter

$$N_a = |\{(x_1, x_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\}\}|$$

$$\text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) = 0,$$

$$\text{Tr}(a_1 x_1 + a_2 x_2) = 0 \}$$

where $a = (a_1, a_2) \in \mathbb{F}_q^2$. By definition and the basic facts of additive characters, for each $a = (a_1, a_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\}$, we have

$$\begin{aligned} N_a &= \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \sum_{y \in \mathbb{F}_p} \chi(y(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1})) \cdot \\ &\quad \sum_{z \in \mathbb{F}_p} \chi(z(a_1 x_1 + a_2 x_2)) - 1 = \\ &= \frac{q^2 - p^2}{p^2} + \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \sum_{y \in \mathbb{F}_p} \chi(y(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1})) + \\ &\quad \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \sum_{z \in \mathbb{F}_p} \chi(z(a_1 x_1 + a_2 x_2)) + \\ &\quad \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \sum_{y, z \in \mathbb{F}_p} \chi(y(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) + \\ &\quad z(a_1 x_1 + a_2 x_2)) =: \\ &= \frac{q^2 - p^2}{p^2} + \Omega_1 + \Omega_2 + \Omega_3 \end{aligned}$$

By Lemmas 1.5 and 1.6, we have

$$\Omega_1 = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{p^{k_1}+1}) \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{p^{k_2}+1}) =$$

$$\begin{cases} \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} G(\eta)\eta(y)(-p^e) & e/d \text{ is odd} \\ \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} G(\eta)\eta(y)(-p^{e+d}) & e/d \text{ is even} \end{cases} = \begin{cases} \frac{p-1}{p^2}(-p^e)G(\eta) & \frac{e}{d} \text{ is odd} \\ \frac{p-1}{p^2}(-p^{e+d})G(\eta) & \frac{e}{d} \text{ is even} \end{cases}$$

By $a = (a_1, a_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\}$, we have

$$\begin{aligned} \Omega_2 &= \frac{1}{p^2} \sum_{z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(za_1 x_1) \cdot \\ &\quad \sum_{x_2 \in \mathbb{F}_q} \chi(za_2 x_2) = 0 \end{aligned}$$

To compute N_a , it is sufficient to determine the value of the exponential sum

$$\begin{aligned} \Omega_3 &= \frac{1}{p^2} \sum_{y, z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{p^{k_1}+1} + za_1 x_1) \cdot \\ &\quad \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{p^{k_2}+1} + za_2 x_2) \end{aligned}$$

For $y \in \mathbb{F}_p^*$, we set $d = \gcd(k, m)$, when m/d is even, $m = 2e$, e/d is even, we know the polynomial $f_1(x) = y^{p^{k_1}} x^{p^{2k_1}} + yx$ is not a permutation polynomial over \mathbb{F}_q ; when m/d is odd or m/d is even, $m = 2e$, e/d is odd, the polynomial $f_1(x) = y^{p^{k_1}} x^{p^{2k_1}} + yx = y(x^{p^{2k_1}} + x)$ must be a permutation polynomial over \mathbb{F}_q .

In fact, suppose that there is $0 \neq b \in \mathbb{F}_q$ such that $f_1(b) = 0$. Then $b^{p^{2k_1}-1} = -1$. Let α be a primitive element of \mathbb{F}_q^* and $b = \alpha^t$, then

$$t(p^{2k_1} - 1) \equiv \frac{p^m - 1}{2} \pmod{p^m - 1} \quad (2)$$

Let $d_1 = \gcd(m, k_1)$, then $\gcd(2k_1, m) = d_1$ by m/d_1 odd. Hence $\gcd(p^{2k_1} - 1, p^m - 1) = (p^{d_1} - 1)$ and $(p^{d_1} - 1) \frac{p^m - 1}{2}$, so Eq. (2) is contradictory.

Since $f_i(x) = y(x^{p^{2k_i}} + x)$ is a permutation polynomial over \mathbb{F}_q , for each $a_i \in \mathbb{F}_q$ there is the unique solution $b_i \in \mathbb{F}_q$ of the equation $x_i^{p^{2k_i}} + x_i + a_i p^{k_i} = 0$. In fact, there is a one-to-one correspondence between $a_i \in \mathbb{F}_q$ and $b_i \in \mathbb{F}_q$, and $a_i = 0$ corresponds to $b_i = 0$.

Hence there is the unique solution $\omega b_i \in \mathbb{F}_q$ of the equation $y(x_i^{p^{2k_i}} + x_i + \omega a_i p^{k_i}) = 0$, where $\omega = \frac{z}{y} \in \mathbb{F}_p^*$.

To compute the value of Ω_3 , we divide into two cases.

The first case: e/d is odd. By Lemma 1.7, we have

$$\begin{aligned} \Omega_3 &= \frac{1}{p^2} \sum_{y, w \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{k_1+1} + ywa_1x_1) \cdot \\ &\quad \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{k_2+1} + ywa_2x_2) = \\ &\quad \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y)G(\eta)(-1)^{\frac{e}{d}} p^e \cdot \\ &\quad \sum_{w \in \mathbb{F}_p^*} \chi'(y \sum_{i=1}^2 (wb_i)^{k_i+1}) = \\ &\quad \frac{-p^e}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y)G(\eta) \cdot \\ &\quad \sum_{w \in \mathbb{F}_p^*} \chi'(yw^2 \text{Tr}(\sum_{i=1}^2 (b_i)^{k_i+1})) \end{aligned}$$

Set

$$\begin{aligned} \Gamma_0 &= \{(b_1, b_2) \in \mathbb{F}_q^2, \{(0, 0)\} \mid \\ &\quad \text{Tr}(\sum_{i=1}^2 b_i^{k_i+1}) = 0\} \end{aligned}$$

$$\Gamma'_0 = \{(b_1, b_2) \in \mathbb{F}_q^2 \mid \text{Tr}(\sum_{i=1}^2 b_i^{k_i+1}) \neq 0\}$$

Since m is even, then we have $\eta(y) = 1$ for $y \in \mathbb{F}_p^*$.

If $(b_1, b_2) \in \Gamma_0$

$$\begin{aligned} \Omega_3 &= \frac{(-p^e)(p-1)^2}{p^2} \cdot G(\eta) \\ N_a &= \frac{q^2 - p^2}{p^2} + \frac{(-p^e)(p-1)}{p^2} G(\eta) + \\ &\quad \frac{(-p^e)(p-1)^2}{p^2} G(\eta) = \\ &\quad \frac{q^2 - p^2}{p^2} + \frac{(-p^e)(p-1)}{p} G(\eta) \end{aligned}$$

Hence, by Lemma 2.1, the weight of C_a is

$$\begin{aligned} n_0 - N_a &= \frac{q^2 - p}{p} + \frac{p-1}{p} (-p^e)G(\eta) - \\ &\quad \frac{q^2 - p^2}{p^2} - \frac{(-p^e)(p-1)}{p} G(\eta) = \\ &\quad p^{2m-1} - p^{2m-2} \end{aligned}$$

If $(b_1, b_2) \in \Gamma'_0$

$$\begin{aligned} \Omega_3 &= \frac{-p^e}{p^2} G(\eta) \sum_{y, w \in \mathbb{F}_p^*} \chi'(yw^2 \text{Tr}(\sum_{i=1}^n b_i^{k_i+1})) = \\ &\quad \frac{(p-1)p^e}{p^2} G(\eta) \\ N_a &= \frac{q^2 - p^2}{p^2} \end{aligned}$$

By Lemma 2.1, the weight of C_a is

$$\begin{aligned} n_0 - N_a &= \frac{q^2 - p}{p} + \frac{p-1}{p} (-p^e)G(\eta) - \frac{q^2 - p^2}{p^2} = \\ &\quad \begin{cases} (p-1)(p^{2m-2} + (-1)^m p^{m-1}) \\ p \equiv 1 \pmod{4} \\ (p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}} p^{m-1}) \\ p \equiv 3 \pmod{4} \end{cases} \end{aligned}$$

The second case: e/d is even. By Lemmas 1.7, 1.8 and 1.9, if $a_2 \in \mathbb{F}_q \setminus S$

$$\begin{aligned} \Omega_3 &= \frac{1}{p^2} \sum_{y, w \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{k_1+1} + ywa_1x_1) \\ &\quad \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{k_2+1} + ywa_2x_2) = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} 0 = 0 \\ N_a &= \frac{q^2 - p^2}{p^2} + \frac{(-p^{e+d})(p-1)}{p^2} G(\eta) \end{aligned}$$

By Lemma 2.1, the weight of C_a is

$$\begin{aligned} n_0 - N_a &= \frac{q^2 - p}{p} + \frac{p-1}{p} (-p^{e+d})G(\eta) - \\ &\quad \frac{q^2 - p^2}{p^2} - \frac{(-p^{e+d})(p-1)}{p^2} G(\eta) = \\ &\quad p^{2m-1} - p^{2m-2} + \frac{(-p^{e+d})(p-1)^2}{p^2} G(\eta) = \\ &\quad \begin{cases} (p-1)(p^{2m-2} + (-1)^m (p-1)p^{m+d-2}) \\ p \equiv 1 \pmod{4} \\ (p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}} (p-1)p^{m+d-2}) \\ p \equiv 3 \pmod{4} \end{cases} \end{aligned}$$

If $a_2 \in S$

$$\begin{aligned} \Omega_3 &= \frac{1}{p^2} \sum_{y, w \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{k_1+1} + ywa_1x_1) \cdot \\ &\quad \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{k_2+1} + ywa_2x_2) = \\ &\quad \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y)G(\eta)(-1)^{\frac{e}{d}+1} p^{e+d} \cdot \\ &\quad \sum_{w \in \mathbb{F}_p^*} \chi'(y \sum_{i=1}^2 (wb_i)^{k_i+1}) = \end{aligned}$$

$$\frac{-p^{e+d}}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y)G(\eta) \sum_{w \in \mathbb{F}_p^*} \chi'(yw^2 \text{Tr}(\sum_{i=1}^2 (b_i)^{k_i+1}))$$

If $(b_1, b_2) \in \Gamma_0$

$$\begin{aligned} \Omega_3 &= \frac{(-p^{e+d})(p-1)^2}{p^2} G(\eta) \\ N_a &= \frac{q^2 - p^2}{p^2} + \frac{(-p^{e+d})(p-1)}{p^2} G(\eta) + \\ &\quad \frac{(-p^{e+d})(p-1)^2}{p^2} G(\eta) = \\ &\quad \frac{q^2 - p^2}{p^2} + \frac{(-p^{e+d})(p-1)}{p} G(\eta) \end{aligned}$$

By Lemma 2.1, the weight of C_a is

$$n_0 - N_a = \frac{q^2 - p}{p} + \frac{p-1}{p}(-p^{\epsilon+d})G(\eta) - \frac{q^2 - p^2}{p^2} - \frac{(-p^{\epsilon+d})(p-1)}{p}G(\eta) = \frac{p^{2m-1} - p^{2m-2}}{p^2}$$

If $(b_1, b_2) \in \Gamma'_0$

$$\Omega_3 = \frac{-p^{\epsilon+d}}{p^2}G(\eta) \sum_{y, w \in \mathbb{F}_p^*} \chi'(y w^2 \text{Tr}(\sum_{i=1}^n b_i^{k_i+1})) = \frac{(p-1)p^{\epsilon+d}}{p^2}G(\eta) N_a = \frac{q^2 - p^2}{p^2}$$

By Lemma 2.1, the weight of C_a is

$$n_0 - N_a = \frac{q^2 - p}{p} + \frac{p-1}{p}(-p^{\epsilon+d})G(\eta) - \frac{q^2 - p^2}{p^2} = \begin{cases} (p-1)(p^{2m-2} + (-1)^m p^{m+d-1}) \\ p \equiv 1 \pmod{4} \\ (p-1)(p^{2m-2} + (-1)^{m+\frac{m}{2}} p^{m+d-1}) \\ p \equiv 3 \pmod{4} \end{cases}$$

Secondly, we determine the frequency of each nonzero weight of C_{D_0} . It is sufficient to consider the values of $|\Gamma_0|, |\Gamma'_0|$.

By Lemma 2.1, it is clear that

$$|\Gamma_0| = n_0 = \begin{cases} \frac{q^2 - p}{p} + \frac{p-1}{p}(-p^{\epsilon})G(\eta) \\ e/d \text{ is odd} \\ \frac{q^2 - p}{p} + \frac{p-1}{p}(-p^{\epsilon+d})G(\eta) \\ e/d \text{ is even} \end{cases}$$

Since $|\Gamma_0| < q^n$. Without loss of generality, suppose that $\Gamma'_0 \neq \emptyset$. If e/d is odd, then for each $c \in \mathbb{F}_p^*$, there are $(x_1, x_2) \in \mathbb{F}_q^2$ such that $\text{Tr}(x_1^{k_1+1} + x_2^{k_2+1}) = c \in \mathbb{F}_p^*$. Hence

$$|\Gamma'_0| = \frac{p-1}{p} \sum_{y \in \mathbb{F}_p} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \chi'(y \text{Tr}(x_1^{k_1+1} + x_2^{k_2+1}) - cy) = \frac{(p-1)q^2}{p} + \frac{p-1}{p} \sum_{y \in \mathbb{F}_p} \chi'(-cy) \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{k_1+1}) \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{k_2+1}) = \frac{(p-1)q^2}{p} + \frac{p-1}{p} \sum_{y \in \mathbb{F}_p} G(\eta) \eta(y) (-p^{\epsilon}) \chi'(-cy) =$$

$$\begin{cases} \frac{p-1}{p}(p^{2m} + (-1)^{m-1} p^m) & p \equiv 1 \pmod{4} \\ \frac{p-1}{p}(p^{2m} + (-1)^{(m-1)+\frac{m}{2}} p^m) & p \equiv 3 \pmod{4} \end{cases}$$

If e/d is even, we have

$$\text{wt}(c_b) \in \{p^{2m-1} - p^{2m-2} + \frac{(-p^{\epsilon+d})(p-1)^2}{p^2}G(\eta) p^{2m-1} - p^{2m-2}, (p-1)(p^{2m-2} + (-1)p^{\epsilon+d-1}G(\eta))\}$$

By Lemma 1.9, $A_{b_1} = p^{2m} - p^{m+2d}$. By the first two Pless Power Moments^[18], the frequency A_{b_i} of b_i satisfies

$$\begin{cases} A_{b_1} + A_{b_2} + A_{b_3} = p^{2m} - 1 \\ b_1 A_{b_1} + b_2 A_{b_2} + b_3 A_{b_3} = p^{2m-1}(p-1)n \end{cases}$$

where $n = \frac{q^2}{p} + \frac{p-1}{p}(-p^{\epsilon+d})G(\eta) - 1$. Hence, we obtain Tables 1, 2.

2.2 The second case

Fix $c \in \mathbb{F}_p^{*2}$ and define

$$D_1 = \{(x_1, x_2) \in \mathbb{F}_q^2 : \text{Tr}(x_1^{k_1+1} + x_2^{k_2+1}) = c\} \\ C_{D_1} = \{c(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2\} \quad (3)$$

where $c(a_1, a_2) = (\text{Tr}(a_1 x_1 + a_2 x_2))_{(x_1, x_2) \in D_1}$.

Since C_{D_1} is a linear code over \mathbb{F}_p , it is independent of the choice of c . For convenience, we take $c=1$.

By Lemma 1.1 and the computation of Γ'_0 as above, we can get the result.

Lemma 2.3 Let $n_1 = |D_1|$. Suppose that e/d is odd, then

$$n_1 = \begin{cases} \frac{1}{p}(p^{2m} + (-1)^{m-1} p^m) \\ p \equiv 1 \pmod{4} \\ \frac{1}{p}(p^{2m} + (-1)^{(m-1)+\frac{m}{2}} p^m) \\ p \equiv 3 \pmod{4} \end{cases}$$

Suppose that e/d is even, then

$$n_1 = \begin{cases} \frac{1}{p}(p^{2m} + (-1)^{m-1} p^{m+d}) \\ p \equiv 1 \pmod{4} \\ \frac{1}{p}(p^{2m} + (-1)^{(m-1)+\frac{m}{2}} p^{m+d}) \\ p \equiv 3 \pmod{4} \end{cases}$$

Theorem 2.4 Let C_{D_1} be the linear code defined as Eq. (3).

If e/d is odd, C_{D_1} is a two-weight code with the Hamming weight distribution in Table 3.

If e/d is even, C_{D_1} is a three-weight code with the Hamming weight distribution in Table 4.

Table 3 e/d is odd

$$p \equiv 1 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)p^{2m-2}$	$\frac{p^{2m}(p+1)}{2p} + \frac{1}{2}(p-1)p^{m-1} - 1$
$(p-1)p^{2m-2} + (-1)^{(m-1)}2p^{m-1}$	$\frac{p-1}{2p}(p^{2m} + (-1)^{m-1}p^m)$

$$p \equiv 3 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)p^{2m-2} + (-1)^{(m-1+\frac{m}{2})}2p^{m-1}$	$\frac{p^{2m}(p+1)}{2p} + \frac{(-1)^{\frac{3m}{2}}}{2} \cdot (p-1)p^{m-1} - 1$
$(p-1)p^{2m-2}$	$\frac{p-1}{2p}(p^{2m} + (-1)^{(m-1)+\frac{m}{2}}p^m)$

Table 4 e/d is even

$$p \equiv 1 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)(p^{2m-2} + (-1)^{m-1}p^{m+d-2})$	$p^{2m} - p^{m+2d}$
$(p-1)p^{2m-2}$	$\frac{p-1}{2}(p^{m+2d-1} - p^{m-d-1})$
$(p-1)p^{2m-2} + (-1)^{(m-1)}2p^{m+d-1}$	$p^{m+2d} - \frac{p-1}{2} \cdot (p^{m+2d-1} - p^{m-d-1}) - 1$

$$p \equiv 3 \pmod{4}$$

Weight	Multiplicity
0	1
$(p-1)(p^{2m-2} + (-1)^{m-1}p^{m+d-2})$	$p^{2m} - p^{m+2d}$
$(p-1)p^{2m-2} + (-1)^{(m-1)}2p^{m+d-1}$	$\frac{p-1}{2}(p^{m+2d-1} - p^{m-d-1})$
$(p-1)p^{2m-2}$	$p^{m+2d} - \frac{p-1}{2} \cdot (p^{m+2d-1} - p^{m-d-1}) - 1$

Proof 3 Firstly, we determine the weight

distribution of the code C_{D_1} . Fix $c=1 \in \mathbb{F}_p^*$. Define the following parameter

$$N'_a = |\{(x_1, x_2) \in \mathbb{F}_q^2 : \text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) = 1, \text{Tr}(a_1x_1 + a_2x_2) = 0\}|$$

where $a = (a_1, a_2) \in \mathbb{F}_q^2 \setminus \{(0, 0)\}$, we have

$$N'_a = \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2, y, z \in \mathbb{F}_p} \chi'(y(\text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) - 1)) \chi(z(a_1x_1 + a_2x_2)) =$$

$$\frac{q^2}{p^2} + \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2, y \in \mathbb{F}_p^*} \chi'(y(\text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) - 1)) \chi(z(a_1x_1 + a_2x_2)) + \frac{1}{p^2}$$

$$+ \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2, z \in \mathbb{F}_p^*} \chi(z(a_1x_1 + a_2x_2)) + \frac{1}{p^2} \sum_{(x_1, x_2) \in \mathbb{F}_q^2, y, z \in \mathbb{F}_p^*} \chi'(y(\text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) - 1)) \cdot \chi(z(a_1x_1 + a_2x_2)) =: \frac{q^2}{p^2} + \Omega'_1 + \Omega'_2 + \Omega'_3$$

We have

$$\Omega'_1 = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \chi'(-y) \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{p^{k_1}+1}) \sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{p^{k_2}+1}) = \begin{cases} -\frac{1}{p^2}G(\eta)(-p^e) & e/d \text{ is odd} \\ -\frac{1}{p^2}G(\eta)(-p^{e+d}) & e/d \text{ is even} \end{cases}$$

By $(a_1, a_2) \in \mathbb{F}_p^2 \setminus \{(0, 0)\}$, we have that

$$\Omega'_2 = 0.$$

Similarly, we have

$$\Omega'_3 = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y)G(\eta)(-p^e) \cdot$$

$$\sum_{w \in \mathbb{F}_p^*} \chi'(y(w^2 \text{Tr}(\sum_{i=1}^2 (b_i)^{p^{k_i}+1}) - 1))$$

where (b_1, b_2) is one-to-one correspondent to (a_1, a_2) .

To compute the value of Ω'_3 , we divide it into two cases.

The first case: e/d is odd.

If $(b_1, b_2) \in \Gamma_0$

$$\Omega'_3 = \frac{p-1}{p^2}G(\eta)(-p^e) \sum_{y \in \mathbb{F}_p^*} \chi'(-y) = (p-1)p^{e-2}G(\eta)$$

$$N'_a = \frac{q^2}{p^2} + p^{e-2}G(\eta) + (p-1)p^{e-2}G(\eta) = p^{2m-2} + p^{e-1}G(\eta) \quad n_1 - N'_a = (p-1)p^{2m-2}$$

If $(b_1, b_2) \in \Gamma'_0$, $w^2 \text{Tr}(\sum_{i=1}^2 (b_i)^{p^{k_i}+1}) \neq 0$ for any $w \in \mathbb{F}_p^*$. By Lemma 1.2

$$\Omega'_3 = \frac{-p^e}{p^2}G(\eta) \left(\sum_{y, w \in \mathbb{F}_p^*} \eta(y) \chi'(y(w^2) \cdot \text{Tr}(\sum_{i=1}^2 (b_i)^{p^{k_i}+1}) - 1) \right) =$$

$$\frac{-p^e}{p^2}G(\eta) \sum_{y \in \mathbb{F}_q^*} \left(\sum_{w \in \mathbb{F}_p} \chi'(y(w^2) \cdot \text{Tr}(\sum_{i=1}^2 (b_i)^{p^{k_i}+1}) - 1) \right) - \chi'(-y) =$$

$$\frac{-p^e}{p^2}G(\eta) - \frac{p^e}{p^2}G(\eta) \sum_{y \in \mathbb{F}_p^*} \eta'(y)$$

$$\begin{aligned} & \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \chi'(-y) G(\eta') = \\ & \frac{-p^e}{p^2} G(\eta) - \frac{-p^e}{p^2} G(\eta) G(\eta')^2 \eta' \cdot \\ & \left(- \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right) \end{aligned}$$

$$N'_a = \frac{q^2}{p^2} + p^{e-2} G(\eta) - p^{e-2} G(\eta) + p^{e-2} \cdot$$

$$G(\eta) G(\eta')^2 \eta' \left(- \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right)$$

$$n_1 - N'_a = \frac{(p-1)q^2}{p^2} + p^{e-1} G(\eta) -$$

$$p^{e-2} G(\eta) G(\eta')^2 \eta' \left(- \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right) =$$

$$\begin{cases} (p-1)p^{2m-2} - p^{m-1} + \\ p^{m-1} \eta' \left(\text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right) & p \equiv 1 \pmod{4} \\ (p-1)p^{2m-2} + (-1)^{\frac{m-1}{2}} p^{m-1} + \\ (-1)^{\frac{m-1}{2}} p^{m-1} \eta' \left(\text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right) & p \equiv 3 \pmod{4} \end{cases}$$

Since $|\Gamma_0| < q^n$. Without loss of generality, suppose that $\Gamma_2 \neq \emptyset$. For some $c \in \mathbb{F}_q^{*2}$, there are $(x_1, x_2) \in \mathbb{F}_q^2$ such that $\text{Tr}(x_1^{p^{k_1}+1} + x_2^{p^{k_2}+1}) = c \in \mathbb{F}_p^{*2}$. By the property of the trace function, the values are presented averagely from \mathbb{F}_p^{*2} .

Hence, set

$$\Gamma_1 = \{ (b_1, b_2) \in \mathbb{F}_q^{*2} \mid \text{Tr} \left(\sum_{i=1}^2 b_i^{p^{k_i+1}} \right) \in \mathbb{F}_p^{*2} \}$$

$$\Gamma_2 = \{ (b_1, b_2) \in \mathbb{F}_q^{*2} \mid \text{Tr} \left(\sum_{i=1}^2 b_i^{p^{k_i+1}} \right) \in \mathbb{F}_p^* \setminus \mathbb{F}_p^{*2} \}$$

$$\begin{aligned} |\Gamma_1| &= \frac{p-1}{2p} \sum_{y \in \mathbb{F}_p} \sum_{(x_1, x_2) \in \mathbb{F}_q^2} \chi'(y \text{Tr}(x_1^{p^{k_1}+1} + \\ & x_2^{p^{k_2}+1}) - cy) = \frac{q^2(p-1)}{2p} + \end{aligned}$$

$$\frac{p-1}{2p} \sum_{y \in \mathbb{F}_p^*} \chi'(-cy) \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{p^{k_1}+1})$$

$$\sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{p^{k_2}+1}) = \frac{(p-1)q^2}{2p} -$$

$$\frac{p-1}{2p} G(\eta) \eta(y) p^e \sum_{y \in \mathbb{F}_p^*} \chi'(-cy) = \frac{(p-1)q^2}{2p} +$$

$$\frac{p-1}{2p} p^e G(\eta) = |\Gamma_2|$$

The second case: e/d is even.

If $a_2 \in \mathbb{F}_q \setminus S$

$$\Omega'_3 = \frac{1}{p^2} \sum_{y, w \in \mathbb{F}_p^*} \chi'(-y) \sum_{x_1 \in \mathbb{F}_q} \chi(yx_1^{p^{k_1}+1} + ywa_1 x_1) \cdot$$

$$\sum_{x_2 \in \mathbb{F}_q} \chi(yx_2^{p^{k_2}+1} + ywa_2 x_2) = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \chi'(-y) 0 = 0$$

$$N'_a = \frac{q^2}{p^2} + p^{e+d-2} G(\eta)$$

By Lemma 2.1, the weight of C_a is

$$n_1 - N'_a = \frac{q^2}{p} + p^{e+d-1} G(\eta) - \frac{q^2}{p^2} - p^{e+d-2} G(\eta) =$$

$$p^{2m-1} - p^{2m-2} + (p-1)p^{e+d-2} G(\eta) =$$

$$\begin{cases} (p-1)(p^{2m-2} + (-1)^{m-1} p^{m+d-2}) \\ p \equiv 1 \pmod{4} \end{cases}$$

$$\begin{cases} (p-1)(p^{2m-2} + (-1)^{m+\frac{m-1}{2}} p^{m+d-2}) \\ p \equiv 3 \pmod{4} \end{cases}$$

If $a_2 \in S$

$$\Omega'_3 = \frac{1}{p^2} \sum_{y \in \mathbb{F}_p^*} \eta(y) G(\eta) (-1)^{\frac{e}{d}+1} p^{e+d} \cdot$$

$$\sum_{w \in \mathbb{F}_p^*} \chi'(y(w^2 \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) - 1))$$

If $(b_1, b_2) \in \Gamma_0$

$$\Omega'_3 = \frac{p-1}{p^2} G(\eta) (-p^{e+d}) \cdot$$

$$\sum_{y \in \mathbb{F}_p^*} \chi'(-y) = (p-1)p^{e+d-2} G(\eta),$$

$$N'_a = \frac{q^2}{p^2} + p^{e+d-1} G(\eta)$$

$$n_1 - N'_a = (p-1)p^{2m-2}$$

If $(b_1, b_2) \in \Gamma'_0$ $w^2 \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \neq 0$ for

any $w \in \mathbb{F}_p^*$. By Lemma 1.2

$$\Omega'_3 = \frac{-p^{e+d}}{p^2} G(\eta) \left(\sum_{y, w \in \mathbb{F}_p^*} \eta(y) \chi'(y(w^2 \cdot$$

$$\text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) - 1) \right) =$$

$$\frac{-p^{e+d}}{p^2} G(\eta) \sum_{y \in \mathbb{F}_p^*} \left(\sum_{w \in \mathbb{F}_p} \chi'(y(w^2$$

$$\text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) - 1) - \chi'(-y) \right) =$$

$$\frac{-p^{e+d}}{p^2} G(\eta) - \frac{-p^{e+d}}{p^2} G(\eta) G(\eta')^2 \eta'$$

$$\left(- \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right)$$

Hence

$$N'_a = \frac{q^2}{p^2} + p^{e+d-2} G(\eta) - p^{e+d-2} G(\eta) +$$

$$p^{e+d-2} G(\eta) G(\eta')^2 \eta' \left(- \text{Tr} \left(\sum_{i=1}^2 (b_i)^{p^{k_i+1}} \right) \right)$$

$$n_1 - N'_a = \frac{(p-1)q^2}{p^2} + p^{e+d-1} G(\eta) - p^{e+d-2} \cdot$$

$$G(\eta)G(\eta')^2\eta'(-\text{Tr}(\sum_{i=1}^2(b_i)^{p^{k_i+1}}))=$$

$$\begin{cases} (p-1)p^{2m-2} - p^{m+d-1} + p^{m+d-1}\eta' \\ (\text{Tr}(\sum_{i=1}^2(b_i)^{p^{k_i+1}})) & p \equiv 1 \pmod{4} \\ (p-1)p^{2m-2} - p^{m+d-1} - p^{m+d-1}\eta' \\ (\text{Tr}(\sum_{i=1}^2(b_i)^{p^{k_i+1}})) & p \equiv 3 \pmod{4} \end{cases}$$

Suppose that

$$b_1 = (p-1)p^{2m-2} + (p-1)p^{e+d-2}G(\eta)$$

$$b_2 = (p-1)p^{2m-2}$$

$$b_3 = (p-1)p^{2m-2} + 2p^{e+d-1}G(\eta)$$

By Lemma 1.9, $A_{b_i} = p^{2m} - p^{m+2d}$. By the first two Pless Power Moments^[18], the frequency A_{b_i} of b_i satisfies the following equations

$$\begin{cases} A_{b_1} + A_{b_2} + A_{b_3} = p^{2m} - 1 \\ b_1 A_{b_1} + b_2 A_{b_2} + b_3 A_{b_3} = p^{2m-1}(p-1)n \end{cases}$$

where $n = \frac{1}{p}(p^{2m} - p^{m+d})$. Hence, we obtain Tables 3,4.

In Eq. (3), suppose that $c \in \mathbb{F}_p^* \setminus \mathbb{F}_p^{*2}$. Similarly, we can get the weight distribution of the linear code. The case can be omitted.

3 Conclusions

There is a recent survey on three-weight codes^[19-23]. We remark that the third class of binary codes is new. We did not find the parameters of the binary three-weight codes of this paper in these literatures.

Linear codes can be used to construct secret sharing schemes^[24]. Let w_{\min} and w_{\max} denote the minimum and maximum nonzero Hamming weights of a linear code C . To obtain secret sharing schemes with interesting access structures, we would like to construct linear codes with the property that

$$w_{\min}/w_{\max} > \frac{p-1}{p}$$

We remark that the linear codes in this paper can be employed in secret sharing schemes using the framework in Ref. [24].

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